

Monte Carlo Sampling Methods in Stochastic Programming

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Graduate Interdisciplinary Program in Applied Mathematics

Oral Comprehensive Exam

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1 Introduction to Stochastic Programming

- Notation and Terminology
- Convenient Special Cases

2 Uncertain Convex Programs

- Uncertain Convex Program
- Theoretical Results
- Computational Results

3 Overlapping Batches

- Bounding Technique for Stochastic Optimization
- Batch Means for Stochastic Programming
- Computational Results

4 Conclusions

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- Convenient Special Cases

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Introduction to Math Programming

$$\begin{aligned} & \min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}) \\ & \text{subject to } \mathbf{x} \in \mathcal{X} \end{aligned}$$

\mathbf{x} Vector of decision variables.

f Objective function.

\mathcal{X} Feasible region.

- e.g., $\mathcal{X} = \{\mathbf{x} \in \mathbb{R}^n : g_i(x) \leq 0, i = 1, \dots, m, \mathbf{x} \geq 0\}$.

Some Terminology

$$\begin{aligned} & \min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}) && \text{(GMP)} \\ & \text{subject to } \mathbf{x} \in \mathcal{X} \end{aligned}$$

- $\hat{\mathbf{x}}$ is **feasible** if $\hat{\mathbf{x}} \in \mathcal{X}$.
- $\hat{\mathbf{x}}$ has **objective function value** or **cost** $f(\hat{\mathbf{x}})$.
- (GMP) is **infeasible** if $\mathcal{X} = \emptyset$.
- \mathbf{x}^* is an **optimal solution** if $\mathbf{x}^* \in \operatorname{argmin}_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x})$.

Theoretically Nice Assumptions

$$\begin{aligned} & \min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}) && \text{(GMP)} \\ & \text{subject to } \mathbf{x} \in \mathcal{X} \end{aligned}$$

Under some conditions

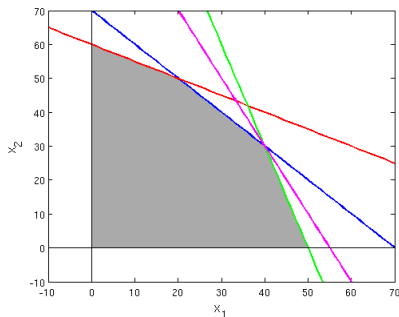
- $f(\mathbf{x})$ is convex in \mathbf{x} .
- \mathcal{X} is a convex set.

Then locally optimal solutions are globally optimal!

- Then (GMP) is called a convex program.

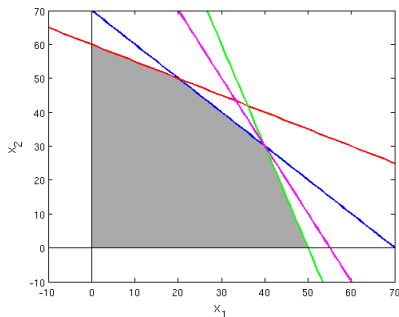
Example: Linear Programming

$$\begin{array}{ll}
 \min & -10x_1 - 8x_2 \\
 \text{subject to} & x_1 + x_2 \leq 70 \\
 & x_1 + 2x_2 \leq 120 \\
 & 3x_1 + x_2 \leq 150 \\
 & 2x_1 + x_2 \leq 110 \\
 & (x_1, x_2) \in \mathbb{R}_+^2
 \end{array}$$



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 \end{array}$$



What if the parameters aren't known?

Introduction to Stochastic Programming

$$\min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}, \xi)$$

subject to $\mathbf{x} \in \mathcal{X}(\xi)$

- \mathbf{x} Vector of decision variables.
- ξ Vector of problem parameters.
- f Objective function, may depend on ξ .
- \mathcal{X} Feasible region, may depend on ξ .
 - e.g.,

$$\mathcal{X}(\xi) = \{\mathbf{x} \in \mathbb{R}^n : g_i(x, \xi) \leq 0, i = 1, \dots, m, \mathbf{x} \geq 0\}.$$

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Uncertain Convex Program

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^n} c^T \mathbf{x} & & (\text{UCP}) \\ \text{subject to } g(\mathbf{x}, \xi) \leq 0, & & \xi \in \Xi \\ \mathbf{x} \in \mathcal{X}. & & \end{aligned}$$

$c^T \mathbf{x}$ Linear objective function, independent of ξ .

$g(\mathbf{x}, \xi)$ Convex scalar constraint, contains all ξ dependence.

\mathcal{X} Feasible region, independent of ξ .

Ξ Support of parameter vector ξ .

Robust Convex Program

- First idea: Find a solution that is feasible for every value of $\xi \in \Xi$.

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^n} c^T \mathbf{x} & \quad (\text{RCP}) \\ \text{subject to } \mathbf{x} \in \mathcal{X} \cap K. \end{aligned}$$

$$K \equiv \bigcap_{\xi \in \Xi} \{\mathbf{x} : g(\mathbf{x}, \xi) \leq 0\}.$$

Analysis of RCP

- One problem: $K = \bigcap_{\xi \in \Xi} \{\mathbf{x} : g(\mathbf{x}, \xi) \leq 0\}$ may be empty.
 - e.g., Ξ is unbounded.
 - Calafiori and Campi assume $K \neq \emptyset$.
- Another problem: (RCP) may be a semi-infinite program.
 - Finitely many variables, infinitely many constraints.
- (RCP) is still a convex program.
 - i.e., $\mathcal{X} \cap K$ is convex.

Chance Constrained Program

$$\min_{\mathbf{x} \in \mathbb{R}^n} c^T \mathbf{x} \quad (\text{CCP})$$

$$\text{subject to } \mathbb{P} \left[g(\mathbf{x}, \tilde{\xi}) > 0 \right] \leq \epsilon$$
$$\mathbf{x} \in \mathcal{X}.$$

\mathbb{P} Probability measure (distribution of $\tilde{\xi}$).

$\tilde{\xi}$ Parameters now take the form of a random variable.

ϵ “Acceptable risk” of infeasible solution.

Analysis of CCP

- Does not suffer from infeasibility
 - i.e., there is (usually) a solution for sufficiently large tolerated risk.
- Feasible region is difficult to calculate.
 - Typically involves multi-dimension integral.
- Feasible region is, in general, **not** convex!

Example of Non-Convex Feasible Region

An example from Sen & Hige:

$$\begin{aligned} & \min_{\mathbf{x} \in \mathbb{R}^2} x_1 + x_2 \\ & \text{subject to } \mathbb{P} \left[\begin{pmatrix} 2 & 1 \\ 1 & 2 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} \geq \begin{pmatrix} \tilde{\xi}_1 \\ \tilde{\xi}_2 \end{pmatrix} \right] \geq 0.5. \end{aligned}$$

Random Variables $\tilde{\xi}_1, \tilde{\xi}_2$ standard Bernoulli distribution.

Dependence $\tilde{\xi}_1 = 1 - \tilde{\xi}_2$.

Feasible Region of Chance Constrained Program

$$\begin{aligned} & \min_{\mathbf{x} \in \mathbb{R}^2} x_1 + x_2 \\ & \text{subject to } \mathbb{P} \left[2x_1 + x_2 \geq \tilde{\xi}_1, x_1 + 2x_2 \geq \tilde{\xi}_2 \right] \geq 0.5. \end{aligned}$$

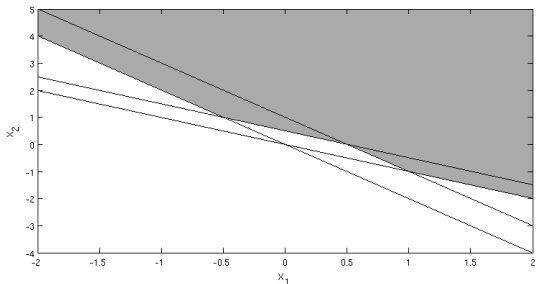


Figure: A graphical representation of feasible region. Feasible region is **not** convex.

Sampled Convex Program

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^n} \quad & c^T \mathbf{x} && (\text{SCP}_N) \\ \text{subject to} \quad & g(\mathbf{x}, \tilde{\xi}^{(i)}) \leq 0, && i = 1, \dots, N \\ & \mathbf{x} \in \mathcal{X}. \end{aligned}$$

$\tilde{\xi}^{(i)}$ Independent and identically distributed (iid) samples of $\tilde{\xi}$.

N Number of samples taken.

Analysis of SCP_N

- Has some danger of infeasibility.
 - Is always feasible if (RCP) is feasible.
- Unlike (RCP), is a finite program!
- Is always a convex program.
- Feasibility of a point is easy to check.
- Solutions of (CCP) and (SCP_N) will outperform those of (RCP).

Benefits of Chance Constrained Problem

$$\min_{x \in \mathbb{R}} x$$

subject to $g(x, \tilde{\xi}) \leq 0$,

$$g(x, \tilde{\xi}) = \begin{cases} \frac{1}{\alpha} - x, & \tilde{\xi} \in [0, \alpha] \\ -x, & \tilde{\xi} \in (\alpha, 1] \end{cases}$$

$$\tilde{\xi} \sim U(0, 1).$$

α A “bad” outcome with small probability.

(RCP) Feasible solutions are $x \geq \frac{1}{\alpha}$.

(CCP) If $\epsilon > \alpha$, feasible solutions are $x \geq 0$.

Optimal Cost $\frac{1}{\alpha}$ or 0.

Relating the Problems

We want to examine solutions to our sampling problem.

- How does the sampled problem compare to the chance constrained problem?
- How does the sampled problem compare to the robust problem?
- How many samples should we take?

Compared to CCP

$$\begin{aligned}
 & \min_{\mathbf{x} \in \mathbb{R}^n} c^T \mathbf{x} && (\text{SCP}_N) \\
 & \text{subject to } g(\mathbf{x}, \tilde{\xi}^{(i)}) \leq 0, && i = 1, \dots, N \\
 & \mathbf{x} \in \mathcal{X}.
 \end{aligned}$$

- Solve (SCP_N) , get \mathbf{x}_N^* .
- What is the probability that \mathbf{x}_N^* violates the constraint?

Principal Result

Theorem (Calafiori & Campi)

Let \mathbf{x}_N^* be the optimal solution to (SCP_N) . Then

$$\mathbb{E}_{\mathbb{P}^N} \left[\mathbb{P}_{\tilde{\xi}} \left[g(\mathbf{x}_N^*, \tilde{\xi}) > 0 \right] \right] \leq \frac{n}{N+1}.$$

n Dimension of problem ($\mathbf{x} \in \mathbb{R}^n$).

\mathbb{P}^N ($= \mathbb{P} \times \dots \times \mathbb{P}$, N times)

- Measure on iid samples $\tilde{\xi}^{(1)}, \dots, \tilde{\xi}^{(N)}$, i.e., on Ξ^N .

$\mathbb{P}_{\tilde{\xi}} = \mathbb{P}$. Independent of \mathbf{x}_N^* .

Discussion of Theorem

$$\mathbb{E}_{\mathbb{P}^N} \left[\mathbb{P}_{\tilde{\xi}} \left[g(\mathbf{x}_N^*, \tilde{\xi}) > 0 \right] \right] \leq \frac{n}{N+1}.$$

- Linear increase with problem size (n) is good.
- Decreases with $1/N$, also good.
- Surprisingly, bound is independent of probability measure!

Moving Beyond Average Behavior

We want to know more than average behavior.

- Use sampling to find \mathbf{x}_N^* .
- Want an estimate on probability of constraint violation.
- What is the probability that our estimate is wrong?
 - ε Level Parameter. Probability of constraint violation is at most ϵ .
 - β Confidence Parameter. Probability that \mathbf{x}_N^* fails to meet desired level parameter.

Principal Result

Corollary (Calafiori & Campi)

Fix two real numbers $\epsilon \in [0, 1]$ (level parameter) and $\beta \in [0, 1]$ (confidence parameter) and let

$$N \geq \frac{n}{\epsilon\beta} - 1. \quad (1)$$

Then

$$\mathbb{P}^N \left[\mathbb{P}_{\tilde{\xi}} \left[g(\mathbf{x}_N^*, \tilde{\xi}) > 0 \right] \leq \epsilon \right] \geq 1 - \beta$$

Discussion of Results

$$N \geq \frac{n}{\epsilon\beta} - 1.$$

- Most surprising: both are independent of the distribution of $\tilde{\xi}$
- Could use when \mathbb{P} is not known explicitly.
 - Provided we can sample from it.
- Linear increase with problem size (n) is good.
- Hyperbolic in $\epsilon\beta$ is not too bad.

Least Squares

- We desire a polynomial of degree $n - 1$ that interpolates the points $(a_i, y_i), i = 1, \dots, m$.
- The problem is

$$\min_{\mathbf{x} \in \mathbb{R}^n} \|A\mathbf{x} - \mathbf{y}\|^2$$

where,

$$A = \begin{bmatrix} 1 & a_1 & \cdots & a_1^{n-1} \\ \vdots & \vdots & & \vdots \\ 1 & a_m & \cdots & a_m^{n-1} \end{bmatrix}, \quad \mathbf{x} = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}, \quad \mathbf{y} = \begin{bmatrix} y_1 \\ \vdots \\ y_m \end{bmatrix}.$$

- Solved by $\mathbf{x} = (A^T A)^{-1} A^T \mathbf{y}$.

Robust Least Squares

- Suppose horizontal coordinates are known inexactly.
 - Belong to some interval
- Matrix A becomes $A(\xi)$.
- Want to solve the problem robustly

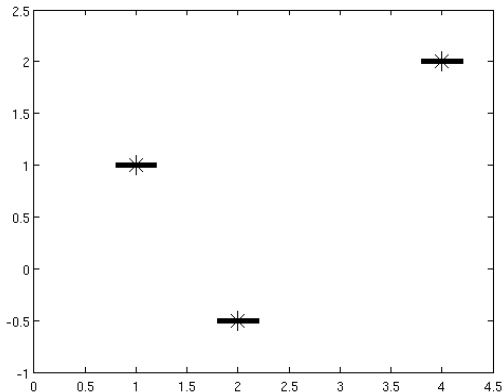
$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\xi \in \Xi} \|A(\xi)\mathbf{x} - \mathbf{y}\|^2$$

- One standard trick:

$$\begin{aligned} & \min_{\mathbf{x} \in \mathbb{R}^n, \gamma \in \mathbb{R}_+} \gamma \\ & \text{subject to } \|A(\xi)\mathbf{x} - \mathbf{y}\|^2 \leq \gamma \quad \forall \xi \in \Xi \end{aligned}$$

Robust Least Squares Example

- Fit a parabola.
- Three (unknown) points
 - $(1, 1)$
 - $(2, -0.5)$
 - $(4, 2)$
- x -coordinate may be ± 0.2 .



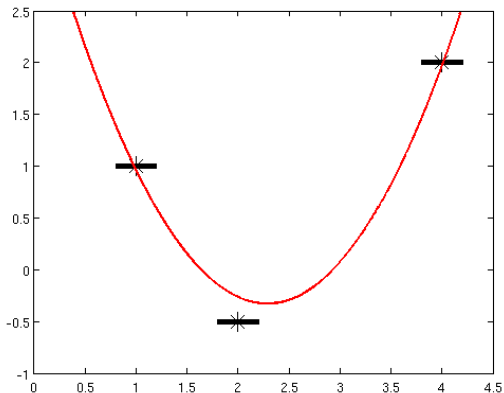
$$\epsilon = 0.1$$

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$$N = 299$$

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- **Solution to Sampled Problem.**



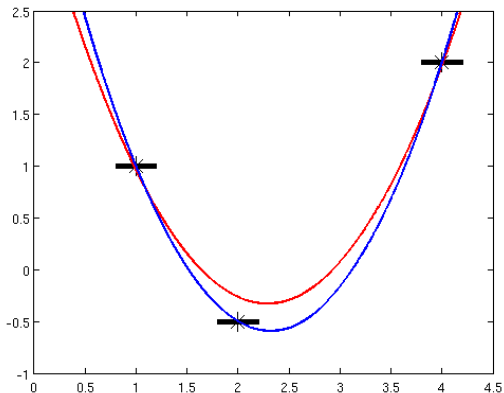
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- **Solution to Sampled Problem.**
- **Expected Value Solution.**



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Estimating Solution Quality

Given a candidate solution, say \mathbf{x}_N^* , we can estimate its probability of infeasibility:

- 1 Generate more iid samples, $\tilde{\xi}^{(1)}, \dots, \tilde{\xi}^{(M')}$.
- 2 Count the number of situations where the constraint is violated. Divide by total number of samples.

$$\hat{V}_{M'}(\mathbf{x}_N^*) = \frac{1}{M'} \sum_{i=1}^{M'} \mathbf{1}_{g(\mathbf{x}_N^*, \tilde{\xi}^{(i)}) > 0}$$

- 3 Then $\hat{V}_{M'}(\mathbf{x}_N^*)$ is an estimate of the constraint violation probability.

Statistics of Solution Quality

$$\mathbb{P}^N \left[\mathbb{P}_{\tilde{\xi}} \left[g(\mathbf{x}_N^*, \tilde{\xi}) > 0 \leq \epsilon \right] \right] \geq 1 - \beta$$

- Hoeffding Inequality allows for statistical analysis of solution quality.
- Want estimate to be accurate to within ϵ' .

$$\mathbb{P}^{M'} \left[\left| \hat{V}_{M'}(\mathbf{x}_N^*) - V(\mathbf{x}_N^*) \right| \leq \epsilon' \right] \geq 1 - 2 \exp\{-2(\epsilon')^2 M'\}$$

- If we desire confidence at least $1 - \beta'$, draw $M' \geq \frac{\log(2/\beta')}{2(\epsilon')^2}$.
- Note: assessing solution quality does not require solving more optimization problems!

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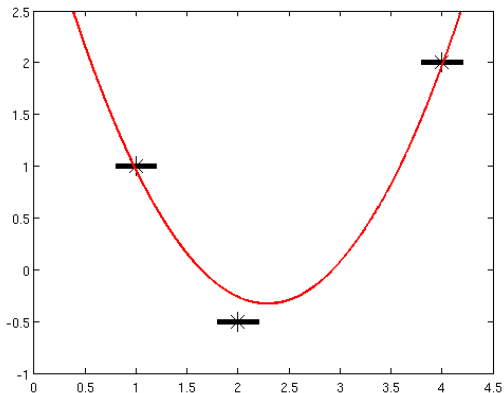
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- **Solution to Sampled Problem.**



$$\epsilon = 0.1$$

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Robust Least Squares Result

- Parameters:

$$\epsilon = 0.1$$

$$\beta = 0.1$$

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- Analysis shows solution has at most 0.6% risk of constraint violation (99.99% confidence).

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Stochastic Objective Function

$$\min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}, \xi)$$

subject to $\mathbf{x} \in \mathcal{X}$

- Now, uncertainty is in the objective function.
- What does it even mean to optimize?
 - Take the “average” cost or benefit?
 - Try to hedge against undesirable outcomes – minimize the probability of something bad?
 - Or maximize the probability of something good?
 - Hedge against undesirable outcomes – Make the “worst case scenario” as good as possible?
- Most general formulation: expected utility...

Stochastic Optimization with Uncertain Cost

$$\begin{aligned} z^* &= \min \mathbb{E}f(\mathbf{x}, \tilde{\xi}) && \text{(SP)} \\ \text{subject to } &\mathbf{x} \in \mathcal{X} \end{aligned}$$

We set up the approximating problem

$$\begin{aligned} z_N^* &= \min \frac{1}{N} \sum_{i=1}^N f(\mathbf{x}, \tilde{\xi}^i) && \text{(SP}_N\text{)} \\ \text{subject to } &\mathbf{x} \in \mathcal{X} \end{aligned}$$

Project Goal

Given a “candidate solution,” $\hat{\mathbf{x}}$, to the stochastic problem, how can we estimate how “good” it is, compared to the optimal solution?

Optimality Gap Estimation Continued

The optimality gap $\mathbb{E}f(\hat{\mathbf{x}}, \tilde{\xi}) - \min_{\mathbf{x} \in \mathcal{X}} \mathbb{E}f(\mathbf{x}, \tilde{\xi})$ is approximated by

$$\frac{1}{M} \sum_{i=1}^M f(\hat{\mathbf{x}}, \tilde{\xi}^i) - \min_{\mathbf{x} \in \mathcal{X}} \frac{1}{M} \sum_{i=1}^M f(\mathbf{x}, \tilde{\xi}^i)$$

- Notice: need to optimize to assess quality!
- To do statistics, we repeat the above multiple times...

Nonoverlapping Batch Means

$N = 12$ Total sample size

$M = 4$ Batch size

$k = \frac{N}{M} = 3$ Number of batches

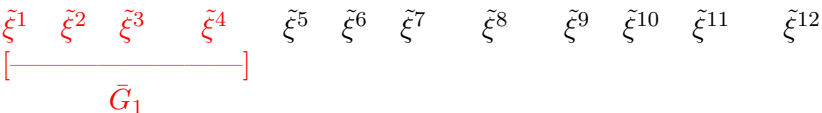
$\tilde{\xi}^1$ $\tilde{\xi}^2$ $\tilde{\xi}^3$ $\tilde{\xi}^4$ $\tilde{\xi}^5$ $\tilde{\xi}^6$ $\tilde{\xi}^7$ $\tilde{\xi}^8$ $\tilde{\xi}^9$ $\tilde{\xi}^{10}$ $\tilde{\xi}^{11}$ $\tilde{\xi}^{12}$

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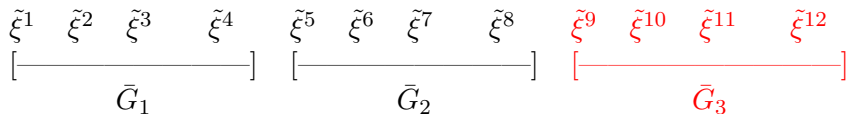
$$\bar{G}_1 = \frac{1}{M} \sum_{i=1}^M f(\hat{\mathbf{x}}, \tilde{\xi}^{M(1-1)+i}) - \min_{\mathbf{x} \in \mathcal{X}} \frac{1}{M} \sum_{i=1}^M f(\mathbf{x}, \tilde{\xi}^{M(1-1)+i})$$

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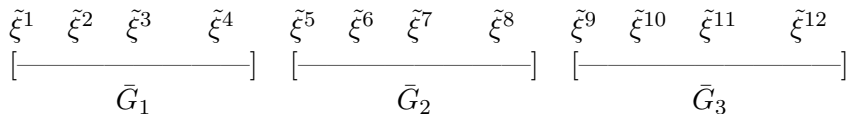
$$\bar{G}_3 = \frac{1}{M} \sum_{i=1}^M f(\hat{\mathbf{x}}, \tilde{\xi}^{M(3-1)+i}) - \min_{\mathbf{x} \in \mathcal{X}} \frac{1}{M} \sum_{i=1}^M f(\mathbf{x}, \tilde{\xi}^{M(3-1)+i})$$

Nonoverlapping Batch Means

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$$\bar{G}_j = \frac{1}{M} \sum_{i=1}^M f(\hat{\mathbf{x}}, \tilde{\xi}^{M(j-1)+i}) - \min_{\mathbf{x} \in \mathcal{X}} \frac{1}{M} \sum_{i=1}^M f(\mathbf{x}, \tilde{\xi}^{M(j-1)+i})$$

$$VG = \frac{1}{(k-1)k} \sum_{j=1}^k (\bar{G}_j - \bar{\bar{G}})^2$$

Batch Means

Why not overlap the batches? (Meketon & Schmeiser, 1984)

Overlapping Batch Means

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$\tilde{\xi}^1$ $\tilde{\xi}^2$ $\tilde{\xi}^3$ $\tilde{\xi}^4$ $\tilde{\xi}^5$ $\tilde{\xi}^6$ $\tilde{\xi}^7$ $\tilde{\xi}^8$ $\tilde{\xi}^9$ $\tilde{\xi}^{10}$ $\tilde{\xi}^{11}$ $\tilde{\xi}^{12}$

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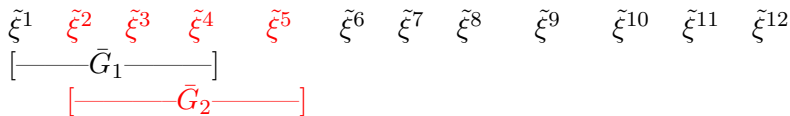
$\tilde{\xi}^1$ $\tilde{\xi}^2$ $\tilde{\xi}^3$ $\tilde{\xi}^4$ $\tilde{\xi}^5$ $\tilde{\xi}^6$ $\tilde{\xi}^7$ $\tilde{\xi}^8$ $\tilde{\xi}^9$ $\tilde{\xi}^{10}$ $\tilde{\xi}^{11}$ $\tilde{\xi}^{12}$
 $[\text{-----}\bar{G}_1\text{-----}]$

Overlapping Batch Means

$N = 12$ Total sample size

$M = 4$ Batch size

$k = \lfloor \frac{N}{M} \rfloor = 3$ Number of nonoverlapping batches

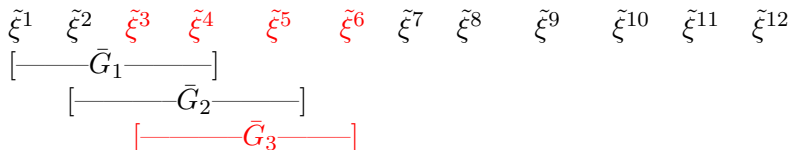


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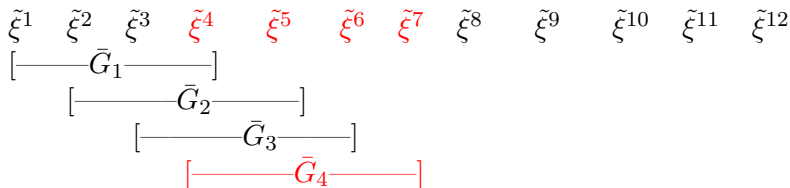


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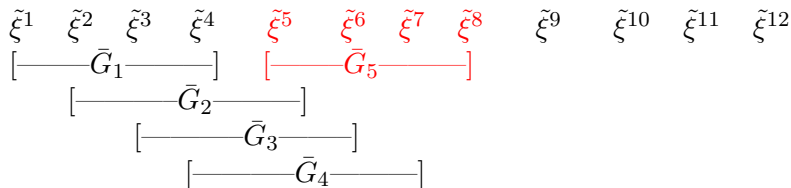


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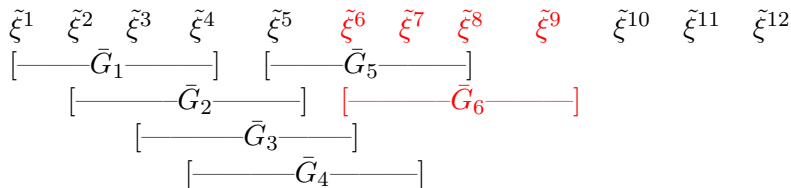


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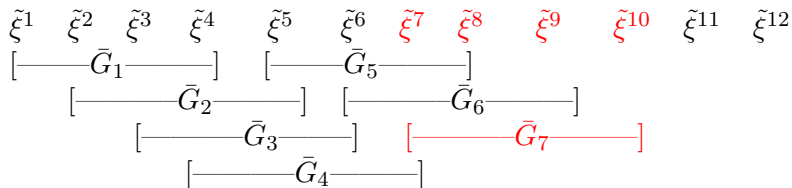


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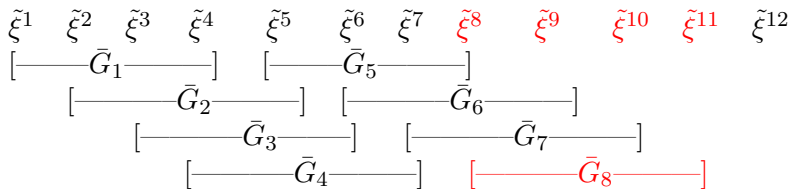


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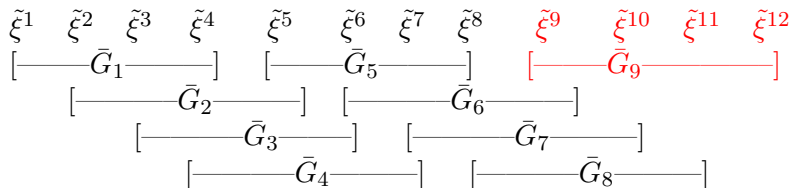


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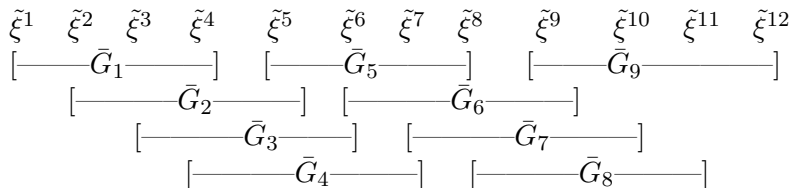


Overlapping Batch Means

$N = 12$ Total sample size

$M = 4$ Batch size

$k = \lfloor \frac{N}{M} \rfloor = 3$ Number of nonoverlapping batches



$$VG = \frac{1}{(k-1)(N-M+1)} \sum_{j=1}^{N-M+1} (\bar{G}_j - \bar{G})^2$$

Partial Overlap

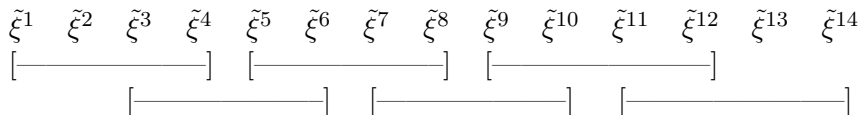


Figure: Graphical representation of $\gamma = 2$.

Let γ be the batch nonoverlap parameter.

- γ gives the number of new samples added to each new batch.
- $\gamma = 1$, maximally overlapping case discussed above.
- $\gamma = M$, classical nonoverlapping case.

Result

- The variance estimator obtained by use of overlapping batches has essentially the same bias as the nonoverlapping estimator.
- Both overlapping and nonoverlapping estimators are unbiased in the limit as $N, M, k \rightarrow \infty$.
- The variance of the sample variance is reduced when overlapping batches is used.

Variance reduction of Sample Variance

- $\gamma \equiv 1$, variance is reduced to $\frac{2}{3}$ of original.
- $\gamma = \frac{1}{3}M$, variance is reduced to $\frac{19}{27}$ of original.
- $\gamma = \frac{1}{2}M$, variance is reduced to $\frac{3}{4}$ of original.
- $\gamma = \frac{1}{L}M$, variance is reduced to $\frac{2L^2 + 1}{3L^2}$ of original.

Theoretical Conclusions

- Generate confidence intervals $\bar{G}_l \pm t_{d,\alpha} \sqrt{V\bar{G}_l}$.
- Confidence intervals are asymptotically valid.

News vendor Problem

- You purchase newspapers from a supplier at cost c , and sell them to the public at price r .
- Number of customers $\tilde{d} \geq 0$ is a random variable.
- Purchase x papers every day.
- Maximize expected profit $\mathbb{E} \left[-cx + r \min \left\{ x, \tilde{d} \right\} \right]$.

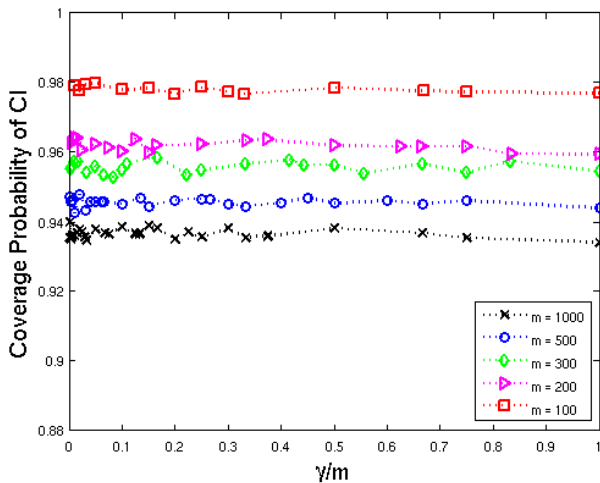
Optimal solution occurs when $\mathbb{P} \left[\tilde{d} \leq x \right] = \frac{r-c}{r}$.

News vendor Problem

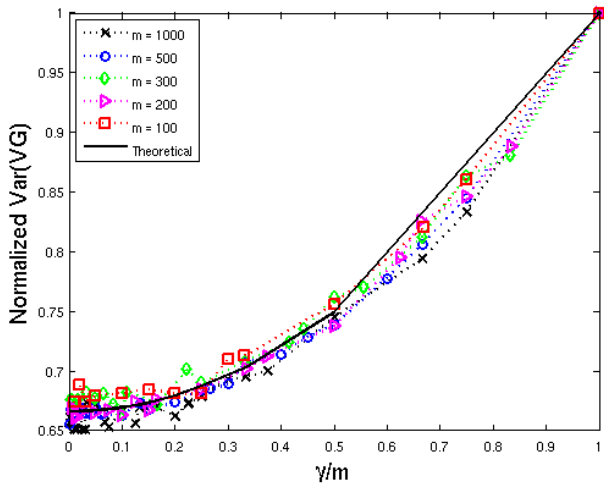
We tested the computational properties of our method on the news vendor problem, with

- Selling price $r = 15$.
- Cost of paper $c = 5$.
- Optimal quantile $\frac{r-c}{r} = \frac{2}{3}$.
- Daily demand of newspapers is $\xi \sim U(0, 10)$.
- $\hat{\mathbf{x}} = 8.7749$.
- To solve:
 - ① Take ξ_1, \dots, ξ_M .
 - ② Sort to get order statistics $\xi_{(1)} \leq \dots \leq \xi_{(M)}$.
 - ③ Optimal solution is $\xi_{(\frac{2}{3}M)}$.

Coverage Probability



Variance Reduction of Variance Estimator



- 1 Introduction to Stochastic Programming
- 2 Uncertain Convex Programs
- 3 Overlapping Batches
- 4 Conclusions

Conclusions

- My work successfully extended results from simulation into stochastic programming.
- By reusing data points, variance of the sample variance is decreased.
- Without sacrificing valid confidence intervals.
- Or increasing bias in sample variance.

Conclusions

- Calafiori and Campi give a new sampling method for solving uncertain convex programs.
- Method is likely to generate solutions that are feasible with high probability.
- Number of samples required is fairly low.
- Assumptions on the problem are fairly general.

Further Work

- Calafiori and Campi's paper inspired much further work.
- I will continue to review this.
- Focusing, in particular, on the “Ambiguous Chance Constrained Program.”
- Learn about computational learning theory, probability metrics.

Questions?

Sketch of Proof

- Let $V(\mathbf{x}) = \mathbb{P} \left[g(\mathbf{x}, \tilde{\xi}) > 0 \right]$.
- Proof proceeds in two steps:
 - Construct an unbiased estimator of $\mathbb{E}_{\mathbb{P}^N} [V(\mathbf{x}^*)]$.
 - Show that the estimator is bounded a.s.

Unbiased Estimator

- Let $\tilde{z}^{(1)}, \dots, \tilde{z}^{(N+1)} \sim \mathbb{P}$ iid.
- Define

$$\min_{\mathbf{x} \in \mathcal{X} \subseteq \mathbb{R}^n} \left\{ c^T \mathbf{x} : g(\mathbf{x}, \tilde{z}^{(i)}) \leq 0, i = 1, \dots, k-1, k+1, \dots, N+1 \right\}, .$$

(SCP_N^k)

- Define indicator variables

$$v_k \equiv \mathbf{1}_{g((\mathbf{x}_N^*)^k, \tilde{z}^{(k)}) > 0}$$

- Then $\hat{V}_N = \frac{1}{N+1} \sum_{k=1}^{N+1} v_k$

Unbiased Estimator Cont.

$$\begin{aligned}\mathbb{E}_{\mathbb{P}^{N+1}} [v_k] &= \mathbb{E}_{\mathbb{P}^N} \left[\mathbb{E}_{\mathbb{P}} \left[v_k | \tilde{z}^{(1)}, \dots, \tilde{z}^{(k-1)}, \tilde{z}^{(k+1)}, \dots, \tilde{z}^{(N+1)} \right] \right] \\ &= \mathbb{E}_{\mathbb{P}^N} \left[\mathbb{P} \left[z \in \Xi : g((\mathbf{x}_N^*)^k, \tilde{z}^{(k)}) > 0 \right] \right] \\ &= \mathbb{E}_{\mathbb{P}^N} \left[V((\mathbf{x}_N^*)^k) \right] \\ &= \mathbb{E}_{\mathbb{P}^N} [V(\mathbf{x}_N^*)],\end{aligned}$$

\hat{V}_N is unbiased.

Estimator is Bounded

Define

$$\min_{\mathbf{x} \in \mathcal{X} \subseteq \mathbb{R}^n} \left\{ c^T \mathbf{x} : g(\mathbf{x}, z^{(i)}) \leq 0, i = 1, \dots, N + 1 \right\}, \quad (\text{SCP}_{N+1})$$

- A constraint is a support constraint if removing it changes the optimal solutions.
- A convex program has, at most, n support constraints.
- Thus solutions to (SCP_N^k) differ from that of (SCP_{N+1}) at most n times.
- Then $\sum_{k=1}^{N+1} v_k \leq n \Rightarrow \hat{V}_N \leq \frac{n}{N+1}$.

◀ Return

Proof of Corollary

- By Markov's Inequality

$$\begin{aligned}\mathbb{P}^N \left[\mathbb{P}_{\tilde{\xi}} \left[g(\mathbf{x}_N^*, \tilde{\xi}) > 0 \right] > \epsilon \right] &\leq \frac{1}{\epsilon} \mathbb{E}_{\mathbb{P}^N} \left[\mathbb{P}_{\tilde{\xi}} \left[g(\mathbf{x}_N^*, \tilde{\xi}) > 0 \right] \right] \\ &\leq \frac{1}{\epsilon} \frac{n}{N+1}\end{aligned}$$

- Set $\frac{1}{\epsilon} \frac{n}{N+1} = \beta$ and solve.

◀ Return

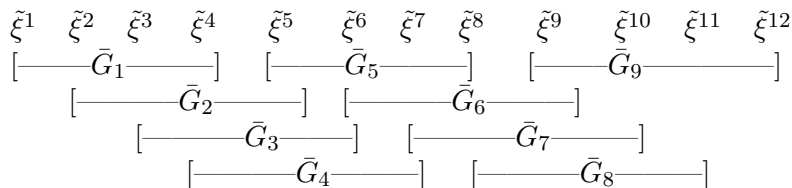
Hoeffding Inequality

- Let $\tilde{\xi}^{(1)}, \dots, \tilde{\xi}^{(M)}$ be independent random variables.
- Let $\mathbb{P} \left[a_i \leq \tilde{\xi}^{(i)} \leq b_i \right] = 1$.
- Let $S = \sum_{i=1}^M \tilde{\xi}^{(i)}$.
- Then the Hoeffding inequality is

$$\mathbb{P} \left[|S - \mathbb{E}[S]| \geq t \right] \leq 2 \exp \left\{ -\frac{2t^2}{\sum_{i=1}^M (b_i - a_i)^2} \right\}$$

- For us, $b_i \equiv 1, a_i \equiv 0, t = M\epsilon$.

Grand Average



- Define the grand mean estimator

$$\bar{G} = \frac{1}{N} \sum_{i=1}^N \frac{1}{|N'(i)|} \sum_{j \in N'(i)} \left[f(\hat{\mathbf{x}}, \tilde{\xi}^i) - f(\mathbf{x}_j^*, \tilde{\xi}^i) \right]$$

◀ Return # 1

◀ Return # 2

Sketch of Proof

$$\bar{G}_j = \frac{1}{M} \sum_{i=1}^M f(\hat{\mathbf{x}}, \tilde{\xi}^{M(j-1)+i}) - \frac{1}{M} \sum_{i=1}^M f(\mathbf{x}_j^*, \tilde{\xi}^{M(j-1)+i})$$

$$\bar{G} = \frac{1}{N} \sum_{i=1}^N \frac{1}{|N'(i)|} \sum_{j \in N'(i)} \left[f(\hat{\mathbf{x}}, \tilde{\xi}^i) - f(\mathbf{x}_j^*, \tilde{\xi}^i) \right]$$

$$\bar{D}_j = \frac{1}{M} \sum_{i=1}^M f(\hat{\mathbf{x}}, \tilde{\xi}^{M(j-1)+i}) - \frac{1}{M} \sum_{i=1}^M f(\mathbf{x}^*, \tilde{\xi}^{M(j-1)+i})$$

$$\bar{D} = \frac{1}{N} \sum_{i=1}^N f(\hat{\mathbf{x}}, \tilde{\xi}^i) - f(\mathbf{x}^*, \tilde{\xi}^i)$$

Sketch Cont.

Show various convergences:

- $\bar{G} - \bar{D} \rightarrow 0$ as $N, M, k \rightarrow \infty$.
- Show $(G_j)^4$ is uniformly integrable.
- $VG - VD \xrightarrow{P} 0$.
- $VG - VD \xrightarrow{L^2} 0$.
- This gives us all the variance reduction results.
- Show that confidence interval estimators $\bar{G} + t_{1-\alpha}\sqrt{VG} - (\bar{D} + t_{1-\alpha}\sqrt{VD}) \xrightarrow{P} 0$.

◀ Return

Optimal Newsvendor Solution

- Suppose \tilde{d} is continuous, with density ϕ , cdf Φ .
- Let $h(x, \tilde{d}) = -cx + r \max\{x, \tilde{d}\}$.

$$\begin{aligned}\mathbb{E} \left[h(x, \tilde{d}) \right] &= -cx + r \mathbb{E} \left[\max\{x, \tilde{d}\} \right] \\ &= -cx + r \int_0^x u \phi(u) du + rx \int_x^\infty \phi(u) du \\ &= -cx + r \int_0^x u \phi(u) du + rx(1 - \Phi(x))\end{aligned}$$

Optimal Newsvendor Solution Cont.

Find the optimal solution by means of the calculus:

$$\begin{aligned}\frac{d}{dx} \mathbb{E} [h(x, \tilde{d})] &= -c + rx\phi(x) + r(1 - \Phi(x)) - rx\phi(x) \\ &= (r - c) - r\Phi(x) \\ &= 0\end{aligned}$$

which gives

$$\Phi(x) = \frac{r - c}{r}$$

◀ Return

Convex Objective Function

$$\begin{aligned} & \min_{\mathbf{x} \in \mathbb{R}^n} f(\mathbf{x}) \\ & \text{subject to } g(\mathbf{x}, \xi) \leq 0, & \xi \in \Xi \\ & \mathbf{x} \in \mathcal{X}. \end{aligned}$$

- New decision variable γ , $\gamma \geq f(\mathbf{x})$.

$$\begin{aligned} & \min_{\mathbf{x} \in \mathbb{R}^n} \gamma \\ & \text{subject to } g(\mathbf{x}, \xi) \leq 0, & \xi \in \Xi \\ & f(\mathbf{x}) - \gamma \leq 0 \\ & \mathbf{x} \in \mathcal{X}. \end{aligned}$$

- Combine with $g'(\mathbf{x}, \xi) = \max\{g(\mathbf{x}, \xi), f(\mathbf{x}) - \gamma\}$.